

EXTENDED CORE AND CHOOSABILITY OF A GRAPH

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ABSTRACT. A graph G is (a, b) -choosable if for any color list of size a associated with each vertices, one can choose a subset of b colors such that adjacent vertices are colored with disjoint color sets. This paper shows an equivalence between the (a, b) -choosability of a graph and the (a, b) -choosability of one of its subgraphs called the extended core. As an application, this result allows to prove the $(5, 2)$ -choosability and $(7, 3)$ -colorability of triangle-free induced subgraphs of the triangular lattice.

1. INTRODUCTION

Let $G = (V(G), E(G))$ be a graph where $V(G)$ is the set of vertices and $E(G)$ is the set of edges, and let a, b, n and e be integers.

Given a list L of G *i.e.* a map $L : V(G) \rightarrow \mathcal{P}(\mathbb{N})$ and a weight ω of G *i.e.* a map $\omega : V(G) \rightarrow \mathbb{N}$, an (L, ω) -choosability c of G is a list of the weighted graph G such that for all $vv' \in E(G)$:

$$c(v) \subset L(v), \quad |c(v)| = \omega(v) \quad \text{and} \quad c(v) \cap c(v') = \emptyset.$$

We say that G is (L, ω) -choosable if there exists an (L, ω) -choosability c of G . An (L, b) -choosability c of G is an (L, ω) -choosability of G such that for all $v \in V(G)$, we have $\omega(v) = b$. A a -list L of G is a list of G such that for all $v \in V(G)$, we have $|L(v)| = a$. The graph G is said to be (a, b) -choosable if for any a -list L of G , there exists an (L, b) -choosability c of G . If the graph is (a, b) -choosable for the a -list L such that $L(v) = L(v')$ for all vertices v, v' , then G is (a, b) -colorable.

The concept of list coloring and choosability was introduced by Vizing [12] and independently by Erdős, Rubin and Taylor [3]. Since then, it has been the subject of many works (see [10, 5, 13, 2, 14] and [6, 8] for more recent papers). In order to characterize 2-choosable (*i.e.* $(2, 1)$ -choosable) graphs, Erdős et al. defined the notion of the core of a graph. The aim of this paper is to extend the notion of core to obtain some characterizations of (a, b) -choosable graphs. Basically, extended cores will be obtained by removing vertices of low degree and induced paths with conditions on the vertex degrees, called handles. Using results on the choosability of a weighted path [1] and extending some of them (Section 2), the main result is Theorem 10 of Section 3 that shows the equivalence between the (a, b) -choosability of a graph G and the (a, b) -choosability of a subgraph of G called its first extended core. Some applications of this theorem for triangle-free

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induced subgraphs of the triangular lattice are given in Section 4 and 5, where it is shown that these graphs are $(5, 2)$ -choosable (Theorem 15) and $(7, 3)$ -colorable (Theorem 21), thus giving another proof of Havet's result [7].

In order to extend a result on the choosability of a weighted path, let us first give some definitions and some known results.

The path P_{n+1} of length n is the graph with vertex set $V = \{v_0, v_1, \dots, v_n\}$ and edge set $E = \bigcup_{i=0}^{n-1} \{v_i v_{i+1}\}$. To simplify the notations $L(i)$ denotes $L(v_i)$ and $c(i)$ denotes $c(v_i)$.

Definition 1. For the path P_{n+1} of length n ,

- a waterfall list L is a list such that for all $i, j \in \{0, \dots, n\}$ with $|i - j| \geq 2$, $L(i) \cap L(j) = \emptyset$;
- two lists L and L' are similar if and only if P_{n+1} is (L, ω) -choosable whenever P_{n+1} is (L', ω) -choosable;
- the amplitude $A_{i,j}(L)$ (or $A_{i,j}$) of a list L is $A_{i,j}(L) = \bigcup_{k=i}^j L(k)$;
- a list L is good if $|L(i)| \geq \omega(i) + \omega(i+1)$ for any $i, 1 \leq i \leq n-1$.

We have proved in [1] the following results:

Proposition 2. For any good list L of P_{n+1} , there exists a similar waterfall list L^c with $|L^c(i)| = |L(i)|$ for all $i \in \{0, \dots, n\}$.

Theorem 3. Let L^c be a waterfall list of a weighted path P_{n+1} . Then P_{n+1} is (L^c, ω) -choosable if and only if:

$$\forall i, j \in \{0, \dots, n\} : |A(i, j)(L^c)| \geq \sum_{k=i}^j \omega(k).$$

Corollary 4. Let L^c be a good waterfall list of a weighted path P_{n+1} such that $|L^c(n)| \geq \omega(n)$. Then P_{n+1} is (L^c, ω) -choosable if and only if

$$\forall j \in \{0, \dots, n\} : |A(0, j)(L^c)| \geq \sum_{k=0}^j \omega(k).$$

Corollary 5. Let L be a list of P_{n+1} such that $|L(0)| = |L(n)| = b$, and $|L(i)| = a = 2b + e$ for all $i \in \{1, \dots, n-1\}$.

If $n \geq \text{Even}\left(\frac{2b}{e}\right)$ then P_{n+1} is (L, b) -choosable

where $\text{Even}(x)$ is the smallest even integer p such that $p \geq x$.

2. CHOOSABILITY OF A PATH

The purpose of this section is to extend the result of Corollary 5 to other constrained lists, namely to 1-reduced lists.

Definition 6. A list L is said to be a 1-reduced list of P_{n+1} if $|L(0)| = b$, for any $i \in \{1, \dots, n-2\} : |L(i)| = a$, $|L(n-1)| = |L(n)| = b + e$, and $|A(n-1, n)(L)| \geq 2b$.

Theorem 7. Let n, a, b, e be four integers such that $a = 2b + e$. Let L be a 1-reduced list of a path P_{n+1} .

If $n = \text{Even}(\frac{2b}{e})$ then P_{n+1} is $(L, b) - \text{choosable}$.

Proof. Since

$$|A(n-1, n)| = |L(n-1)| + |L(n) \setminus L(n-1)| = (b+e) + |L(n) \setminus L(n-1)| \geq 2b,$$

we obtain that $|L(n) \setminus L(n-1)| \geq b - e$. Let D be a set such that

$$D \subset L(n) \setminus L(n-1) \text{ and } |D| = b - e.$$

Let L' be the new list constructed with L such that (see the following figure):

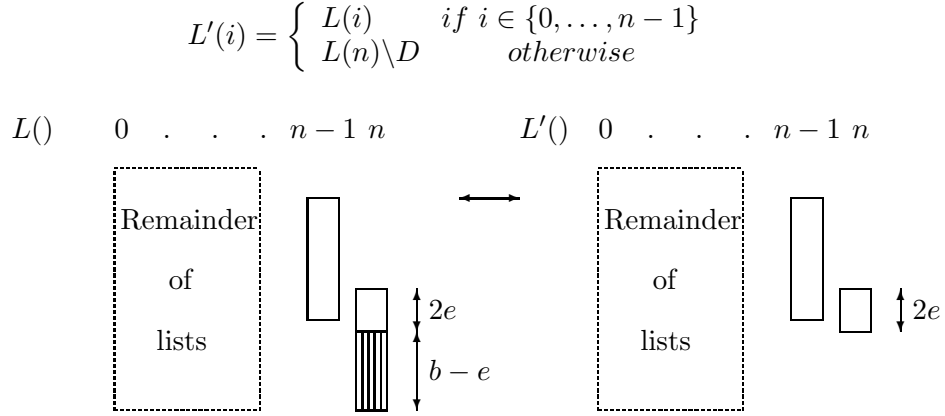


Fig. 1. Construction of L' (on the right) for a 1-reduced list L (on the left).

A new weight function ω' is constructed such that:

$$\omega'(i) = \begin{cases} b & \text{if } i \in \{0, \dots, n-1\} \\ e & \text{otherwise.} \end{cases}$$

We are going to prove that:

if P_{n+1} is $(L', \omega') - \text{choosable}$ then P_{n+1} is $(L, b) - \text{choosable}$.

Indeed, if c' is an (L', ω') -choosability of P_{n+1} , then we construct c such that:

$$c(i) = \begin{cases} c'(i) & \text{if } i \in \{0, \dots, n-1\} \\ c'(n) \cup D & \text{otherwise.} \end{cases}$$

Since $D \cap L(n-1) = \emptyset$, we have $c(n-1) \cap c(n) = \emptyset$ and then c is an (L, b) -choosability of P_{n+1} .

Now, this new list L' is a good list of P_{n+1} and $|L'(n)| \geq \omega'(n)$. Proposition 2 shows that there exists a waterfall list L^c similar to L' such that for all k we have $|L^c(k)| = |L'(k)|$.

Thanks to Corollary 4, it remains to check that:

$$\forall j \in \{0, \dots, n\} : |A(0, j)(L^c)| \geq \sum_{k=0}^j \omega'(k).$$

Case 1: $j \in \{0, \dots, n-2\}$. Since the list is a waterfall list, we have:

$$|A(0, j)(L^c)| \geq \begin{cases} \sum_{\substack{k=0 \\ k \text{ even}}}^j |L^c(k)| = b + a\frac{j}{2} & \text{if } j \text{ is even} \\ \sum_{\substack{k=0 \\ k \text{ odd}}}^j |L^c(k)| = a\frac{j+1}{2} & \text{otherwise} \end{cases}$$

and the weight function satisfies $\sum_{k=0}^j \omega'(k) = (j+1)b$. Hence, we deduce that

$$|A(0, j)(L^c)| \geq \sum_{k=0}^j \omega'(k).$$

Case 2: $j = n-1$. Since the list is a waterfall list, we have:

$$|A(0, n-1)(L^c)| \geq \sum_{\substack{k=0 \\ k \text{ odd}}}^{n-1} |L^c(k)| = (b+e) + a\frac{n-2}{2},$$

and $\sum_{k=0}^{n-1} \omega'(k) = nb$. Then $b+e+a\frac{n-2}{2} \geq nb$ if and only if $\frac{en}{2} \geq b$, which is true by hypothesis since $n = \text{Even}(\frac{2b}{e})$, thus

$$|A(0, n-1)(L^c)| \geq \sum_{k=0}^{n-1} \omega'(k).$$

Case 3: $j = n$. Since the list is a waterfall list, we have:

$$|A(0, n)(L^c)| \geq \sum_{\substack{k=0 \\ k \text{ even}}}^n |L^c(k)| = b + 2e + a\frac{n-2}{2},$$

and $\sum_{k=0}^n \omega'(k) = nb + e$, then $b + 2e + a\frac{n-2}{2} \geq nb + e$ if and only if $\frac{en}{2} \geq b$, which is true by hypothesis since $n = \text{Even}(\frac{2b}{e})$, thus

$$|A(0, n)(L^c)| \geq \sum_{k=0}^n \omega'(k).$$

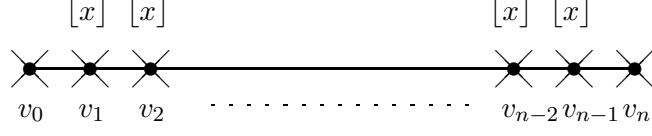
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3. THE EXTENDED CORE OF A GRAPH

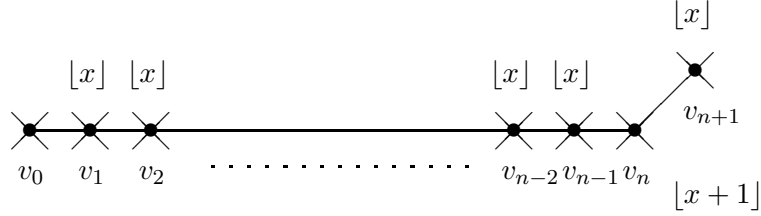
The purpose of this section is to prove Theorem 10 which gives the equivalence between the (a, b) -choosability of a graph G and the (a, b) -choosability of one of its subgraphs, its first extended core, denoted $\text{Core}_{ch}(x, 1)(G)$ (with the idea to take $x = \frac{a}{b}$).

The first extended core is a generalization of the core, introduced by Erdős, Rubin and Taylor in [3] for 2-choosable graphs.

Definition 8. A handle_x of length n in a graph G is a path $\{v_0, \dots, v_n\}$ such that the vertices $\{v_1, \dots, v_{n-1}\}$ have degree less or equal to $\lfloor x \rfloor$ and for all $i, j \in \{1, \dots, n-1\}$ we have: if $|i-j| \geq 2$ then $v_i v_j \notin E$. The interior of the handle_x is the set of vertices $\{v_1, \dots, v_{n-1}\}$.

Fig. 2. Example of a handle_x .

Definition 9. A 1-handle_x of length n is a handle_x of length n such that v_n has degree less or equal to $\lfloor x+1 \rfloor$ and has a neighbor v_{n+1} of degree less or equal to $\lfloor x \rfloor$ and for all $i, j \in \{1, \dots, n+1\}$ we have: if $|i-j| \geq 2$ then $v_i v_j \notin E$.

Fig. 3. Example of a 1-handle_x .

We define the first extended core, denoted $\text{Core}_{ch}(x, 1)(G)$, of a graph G as the induced subgraph of G obtained inductively when we remove:

- its vertices of degree $0, 1, \dots, \lfloor x-1 \rfloor$,
- the interior of its handle_x of length $n \geq \text{Even}(\frac{2}{x-\lfloor x \rfloor})$,
- the interior of its 1-handle_x of length $n = \text{Even}(\frac{2}{x-\lfloor x \rfloor}) - 1$.

Let us remark that the definition of $\text{Core}_{ch}(x, 1)(G)$ does not depend on the order which we use to remove the vertices and that the subscript 'ch' means that we will use it for choosability purpose.

Theorem 10. Let a, b be two integers and x be a rational number such that $\frac{a}{b} \geq x$. For any graph G , we have the following equivalence:

$$G \text{ is } (a, b) \text{ -- choosable} \Leftrightarrow \text{Core}_{ch}(x, 1)(G) \text{ is } (a, b) \text{ -- choosable}.$$

Proof. Since $\text{Core}_{ch}(x, 1)(G)$ is a subgraph of G , if G is (a, b) -choosable then $\text{Core}_{ch}(x, 1)(G)$ is of course (a, b) -choosable.

Conversely, suppose that $\text{Core}_{ch}(x, 1)(G)$ is (a, b) -choosable. Let us discuss the different cases and let e be the integer defined by:

$$e = a - \lfloor x \rfloor b \geq 0.$$

Case 1: we remove a vertex of low degree. Let $v \in V(G)$ such that its degree $d(v) \leq \lfloor x \rfloor - 1$. Let $N(v)$ be the set of neighbors of v in G . Suppose that $G - \{v\}$ is (a, b) -choosable and let L be an a -list of G . Then there exists an (L, b) -choosability c of $G - \{v\}$.

Let $L'(v) = L(v) \setminus \bigcup_{w \in N(v)} c(w)$. Then

$$|L'(v)| \geq |L(v)| - d(v)b \geq a - (\lfloor x \rfloor - 1)b = b + e \geq b.$$

Hence, we can complete the choosability with $c(v) \subset L'(v)$ such that $|c(v)| = b$, and thus G is (a, b) -choosable.

Case 2: x is an integer ($x = \lfloor x \rfloor$). Since $\text{Even}(\frac{2}{x-\lfloor x \rfloor}) = \infty$, the extended core $\text{Core}_{ch}(x, 1)(G)$ is limited to the case 1, and the proof is done.

Case 3: x is not an integer ($x > \lfloor x \rfloor$). It remains two kinds of handle_x to consider:

Subcase 1: Let $H_n(x)$ be a handle_x of G of length $n \geq \text{Even}(\frac{2}{x-\lfloor x \rfloor})$. Suppose that $G' = G - \{v_1, \dots, v_{n-1}\}$ is (a, b) -choosable, and let L be an a -list of G . Then there exists an (L, b) -choosability c of G' . We set $N'(v_i)$ to be the set of neighbors of v_i in the induced subgraph $G - \{v_{i-1}, v_{i+1}\}$. By hypothesis for all $i \in \{1, \dots, n-1\}$, the degree of v_i satisfies $d(v_i) \leq \lfloor x \rfloor$ hence $|N'(v_i)| \leq \lfloor x \rfloor - 2$. Let L' be the list of $H_n(x)$ such that:

$$L'(i) = \begin{cases} c(v_i) & \text{if } i \in \{0, n\} \\ L(v_i) \setminus \bigcup_{w \in N'(v_i)} c(w) & \text{otherwise.} \end{cases}$$

Then for all $i \in \{1, \dots, n-1\}$, we have:

$$|L'(i)| \geq |L(v_i)| - \sum_{w \in N'(v_i)} |c(w)| \geq a - (\lfloor x \rfloor - 2)b = 2b + e.$$

Suppose without loss of generality that for all $i \in \{1, \dots, n-1\}$, we have: $|L'(i)| = 2b + e$. Since $\frac{a}{b} \geq x$ we have $\frac{2}{x-\lfloor x \rfloor} \geq \frac{2b}{e}$ and thus

$$n \geq \text{Even}(\frac{2}{x-\lfloor x \rfloor}) \geq \text{Even}(\frac{2b}{e}).$$

Since $|L'(0)| = |L'(n)| = b$, then L' is a list which satisfies the hypothesis of Corollary 5: we obtain the existence of an (L', b) -choosability c' of P_{n+1} , *i.e.* of $H_n(x)$. Finally, we construct an (L, b) -choosability c'' of G such that:

$$c''(v) = \begin{cases} c(v) & \text{if } v \in G' \\ c'(v) & \text{otherwise.} \end{cases}$$

Subcase 2: Let $H_n(x)$ be a 1-handle $_x$ of G of length $n = \text{Even}(\frac{2}{x-\lfloor x \rfloor}) - 1$. If $\text{Even}(\frac{2}{x-\lfloor x \rfloor}) > \text{Even}(\frac{2b}{e})$ then $n \geq \text{Even}(\frac{2b}{e})$, it is a handle_x of sufficiently big length, and hence we come back to Subcase 1. Otherwise, $\text{Even}(\frac{2}{x-\lfloor x \rfloor}) = \text{Even}(\frac{2b}{e})$ (because $\text{Even}(\frac{2}{x-\lfloor x \rfloor}) \geq \text{Even}(\frac{2b}{e})$). By definition, there exists a vertex $v_{n+1} \in G - H_n(x)$ which is a neighbor of v_n such that $d(v_{n+1}) \leq \lfloor x \rfloor$. Suppose that $G' = G - \{v_1, \dots, v_{n-1}\}$ is (a, b) -choosable, and let L be an a -list of G and c be an (L, b) -choosability of G' . For any $i \in \{1, \dots, n\}$, $N'(v_i)$ is the set of neighbors of v_i in $G - \{v_{i-1}, v_{i+1}\}$ and $N'(v_{n+1})$ is the set of neighbors of v_{n+1} in $G - \{v_n\}$. Let L' be the list of $H_n(x) \cup \{v_{n+1}\}$ such that:

$$L'(i) = \begin{cases} c(v_0) & \text{if } i = 0 \\ L(v_i) \setminus \bigcup_{w \in N'(v_i)} c(w) & \text{otherwise.} \end{cases}$$

For any $i \in \{1, \dots, n-1\}$, we have:

$$|L'(i)| \geq |L(v_i)| - \sum_{w \in N'(v_i)} |c(w)| \geq a - (\lfloor x \rfloor - 2)b = 2b + e.$$

Furthermore, for $i \in \{n, n+1\}$:

$$|L'(i)| \geq |L(v_i)| - \sum_{w \in N'(v_i)} |c(w)| \geq a - (\lfloor x \rfloor - 1)b = b + e.$$

We can suppose without loss of generality that for all $i \in \{1, \dots, n-1\}$, we have $|L'(i)| = 2b + e$, and for $i \in \{n, n+1\}$, we have $|L'(i)| = b + e$. Since c is an (L, b) -choosability of G' then $c(v_n) \cap c(v_{n+1}) = \emptyset$ and since for $i \in \{n, n+1\}$, we have $c(v_i) \subset L'(i)$ then we obtain:

$$|A(n, n+1)(L')| = |L'(n) \cup L'(n+1)| \geq |c(v_n)| + |c(v_{n+1})| = 2b.$$

Hence L' is a 1-reduced list of $H_n(x) \cup \{v_{n+1}\}$; but since $n+1 = \text{Even}(\frac{2b}{e})$, Theorem 7 constructs an (L', b) -choosability c' of $H_n(x) \cup \{v_{n+1}\}$. Finally, we construct an (L, b) -choosability c'' of G such that:

$$c''(w) = \begin{cases} c(w) & \text{if } w \in G - (H_n(x) \cup \{v_{n+1}\}) \\ c'(w) & \text{otherwise.} \end{cases}$$

In conclusion, if $\text{Core}_{ch}(x, 1)(G)$ is (a, b) -choosable, we add successively the subgraphs of G that we removed (to obtain the $\text{Core}_{ch}(x, 1)$), in the opposite order. Thanks to the cases studied in every step, the graph remains (a, b) -choosable and thus G is (a, b) -choosable. \square

Let us define the set $\mathcal{C}_h(x)$ to be the set of graphs G which are (a, b) -choosable for all a, b such that $\frac{a}{b} \geq x$, i.e.:

$$\mathcal{C}_h(x) = \{G, \text{ such that for all } \frac{a}{b} \geq x, G \text{ is } (a, b)\text{-choosable}\}.$$

Hence, we deduce the following corollary:

Corollary 11. *Let G be a graph and x be a rational number. Then:*

$$G \in \mathcal{C}_h(x) \Leftrightarrow \text{Core}_{ch}(x, 1)(G) \in \mathcal{C}_h(x).$$

4. APPLICATION TO THE TRIANGULAR LATTICE

Let \mathcal{R} be a triangle-free induced subgraph of the triangular lattice. Recall that the triangular lattice is embedded in an Euclidian space and that any vertex (x, y) of \mathcal{R} has at most six neighbors: its neighbor on the left, $(x-1, y)$, its neighbor on the right $(x+1, y)$, its neighbor on the top left $(x-1, y+1)$, its neighbor on the top right $(x, y+1)$, its neighbor on the bottom left $(x, y-1)$ and its neighbor on the bottom right $(x+1, y-1)$.

Definition 12. *The nodes of \mathcal{R} are the vertices of degree 3. There are two kinds of nodes: the left nodes whose neighbors are the neighbors on the left, on the top right, and on the bottom right; and the right nodes whose neighbors are the neighbors on the right, on the top left and on the bottom left.*



Fig. 4. Left node and right node.

Definition 13. A cutting node of \mathcal{R} is a left node (x, y) such that for any node (x', y') , we have $y \geq y'$, and for any left node (x', y) , we have $x' \leq x$. A handle of \mathcal{R} is a path P such that its extremal vertices are nodes and its internal vertices have degree 2.

A cutting handle of \mathcal{R} is a handle such that one of its extremal vertices is the cutting node (x, y) and one of its internal vertices is $(x, y + 1)$.

One can see the cutting node as the left node the most on the top on the right.

We have the trivial following lemma:

Lemma 14. Let P_{n+1} be a cutting handle of \mathcal{R} such that $V(P_{n+1}) = \{v_0, \dots, v_n\}$. If the length n of P_{n+1} is less or equal to 3, then $n = 3$ and v_3 has got a neighbor $v_4 \neq v_2$ of degree less or equal to 2.

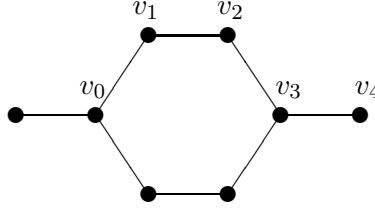


Fig. 5. Cutting handle of length 3.

We have:

Theorem 15. If \mathcal{R} is a triangle-free induced subgraph of the triangular lattice, then:

$$\mathcal{R} \in \mathcal{C}_h\left(\frac{5}{2}\right).$$

Proof. We set $x = \frac{5}{2}$, thus $\lfloor x \rfloor = 2$, hence a handle of \mathcal{R} is a handle $_x$ of \mathcal{R} . Let $G = \text{Core}_{ch}(x, 1)(\mathcal{R})$. If $G = \emptyset$ then Corollary 11 gives the result. Otherwise $G \neq \emptyset$, and since the girth of this graph is at least 6 (and since we have removed all the handle $_x$ of length $n \geq \text{Even}(\frac{2}{x - \lfloor x \rfloor}) = 4$), G can't be a forest of cycles and thus G has got at least two nodes (since we have removed the vertices of degree 0 and 1, the number of nodes is necessarily even). By symmetry, one can suppose that G has a cutting handle (otherwise we can consider its mirror graph). Since $x = \frac{5}{2}$ and since we have removed all the handle $_x$ of length $n \geq 4$, then Lemma 14 shows that this cutting handle P_{n+1} has length $n = 3$, $V(P_3) = \{v_0, \dots, v_3\}$ and v_3 has a neighbor $v_4 \neq v_2$ of degree less or equal to 2. Thus P_3 is a 1-handle $_x$, but this is absurd since $G = \text{Core}_{ch}(x, 1)(\mathcal{R})$. \square

5. THE SECOND EXTENDED CORE

We can define a second extended core by removing more vertices.

Definition 16. A 2-handle_x of length n in a graph G is a 1-handle_x of length n such that v_0 has degree less or equal to $\lfloor x+1 \rfloor$ and has a neighbor v_{-1} of degree less or equal to $\lfloor x \rfloor$ and for all $i, j \in \{-1, \dots, n+1\}$ we have: if $|i - j| \geq 2$ then $v_i v_j \notin E$.

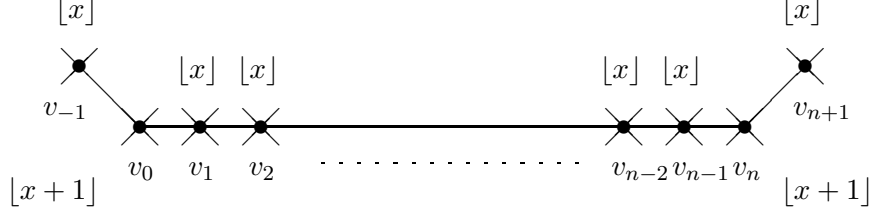


Fig. 6. Example of a 2-handle_x .

Then, we define the second extended core, denoted $\text{Core}_{ch}(x, 2)(G)$, of a graph G as the induced subgraph of G obtained inductively when we remove:

- its vertices of degree $0, 1, \dots, \lfloor x-1 \rfloor$,
- the interior of its handle_x of length $n \geq \text{Even}(\frac{2}{x-\lfloor x \rfloor})$,
- the interior of its 1-handle_x of length $n = \text{Even}(\frac{2}{x-\lfloor x \rfloor}) - 1$,
- the interior of its 2-handle_x of length $n = \text{Even}(\frac{2}{x-\lfloor x \rfloor}) - 2$.

We can prove (see [4] for the details) in the same way as for the first extended core that the choosability of a graph reduces to the choosability of its second extended core:

Theorem 17. Let a, b be two integers and x be a rational number such that $\frac{a}{b} \geq x$. For any graph G , we have the following equivalence:

$$G \text{ is } (a, b) - \text{choosable} \Leftrightarrow \text{Core}_{ch}(x, 2)(G) \text{ is } (a, b) - \text{choosable}.$$

Definition 18. Let $x \in [2, 3[$ be a rational number. A handle_x of parity of a graph G is a handle_x of length n , such that there exists another path in G of length m , from v_0 to v_n , with $n \geq m$ and $m \equiv n \pmod{2}$.

Definition 19. The $\text{Core}_{co}(x, 2)(G)$ is defined as the $\text{Core}_{ch}(x, 2)(G)$, but in addition, we remove also, for $x \in [2, 3[$, the interior of its handle_x of parity.

Let us define the set $\mathcal{C}_o(x)$ to be the set of graphs G which are (a, b) -colorable for all a, b such that $\frac{a}{b} \geq x$, i.e.:

$$\mathcal{C}_o(x) = \{G, \text{ such that for all } \frac{a}{b} \geq x, G \text{ is } (a, b)\text{-colorable}\}.$$

The McDiarmid and Reed conjecture (see [9]) asserts that:

$$\mathcal{R} \in \mathcal{C}_o(\frac{9}{4}).$$

One can prove in the same way as for the $\text{Core}_{ch}(x, 2)(G)$ (see [4] for the details) the following theorem:

Theorem 20. Let $x \in [2, 3[$. For any graph G , we have:

$$G \in \mathcal{C}_o(x) \Leftrightarrow \text{Core}_{co}(x, 2)(G) \in \mathcal{C}_o(x).$$

Then, Theorem 20 enables us to find again a result of Havet on the coloration of the graph \mathcal{R} :

Theorem 21 (Havet, [7]).

$$\mathcal{R} \in \mathcal{C}_o(\frac{7}{3}) .$$

Proof. We set $x = \frac{7}{3}$, and $G = \text{Core}_{co}(x, 2)(\mathcal{R})$. If $G = \emptyset$ then Theorem 20 gives the proof. Otherwise, $G \neq \emptyset$ and there exists a cutting handle P_{n+1} of G . Since $\text{Even}(\frac{2}{x-\lfloor x \rfloor}) = 6$, we have removed all the handle $_x$ of length $n \geq 6$. For $n \leq 5$, a careful study (see [4] where some technical lemmas are used) shows that we obtain a contradiction, which concludes the proof. \square

Remark 22. Using the $\text{Core}_{ch}(7/3, 2)$, we can almost prove (some few cases still resist) that $\mathcal{R} \in \mathcal{C}_h(\frac{7}{3})$.

Remark 23. Recall that a graph G is said to be (a, b) -free-choosable in a vertex v_0 if for any list L of G such that for any $v \in V(G) \setminus \{v_0\} : |L(v)| = a$ and $|L(v_0)| = b$, there exists an (L, b) -choosability c of G . If G is a graph such that $\text{Core}_{ch}(x, 2)(G) = \emptyset$, then we have shown in particular that if a, b are such that $\frac{a}{b} \geq x$ then G is (a, b) -free-choosable in the last vertex removed by the algorithm used to get the core. In particular, Theorem 15 implies that there exists a vertex v_0 for which the graph \mathcal{R} is $(5, 2)$ -free-choosable in v_0 .

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